

# LTE-A Spectral Resource Management

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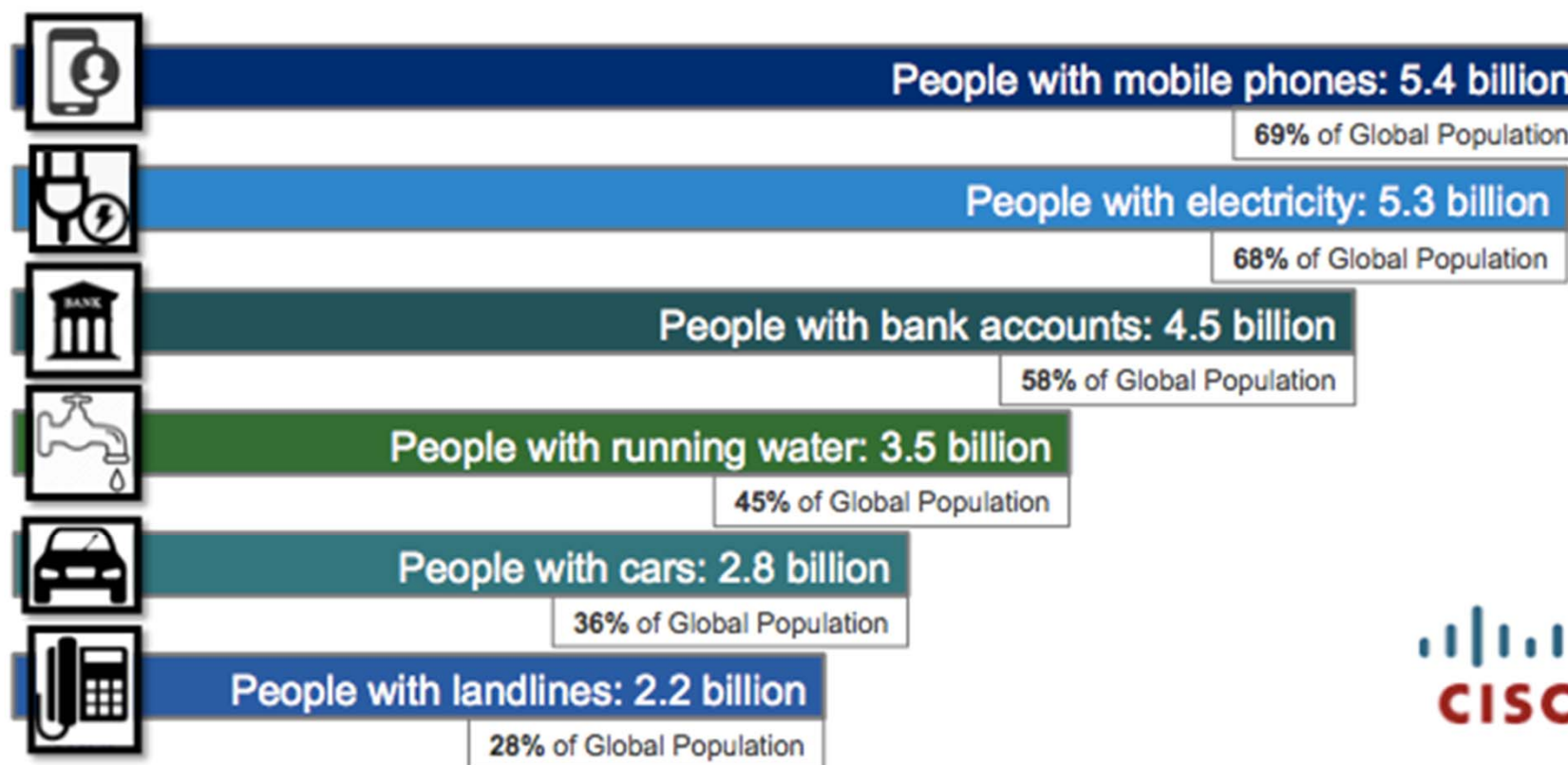
# Outline

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- Brief introduction to LTE-A
- A 4-guideline downlink scheduling strategy
- Downlink resource management for green communication

# Cisco forecast (1/2)

- By 2020:



# Cisco forecast (2/2)

**By 2019:**



## IP Broadband Growth Projections

More Internet  
Users



| 2014        | 2019        |
|-------------|-------------|
| 2.8 Billion | 3.9 Billion |

More Devices &  
Connections



| 2014         | 2019         |
|--------------|--------------|
| 14.2 Billion | 24.4 Billion |

Faster Broadband  
Speeds



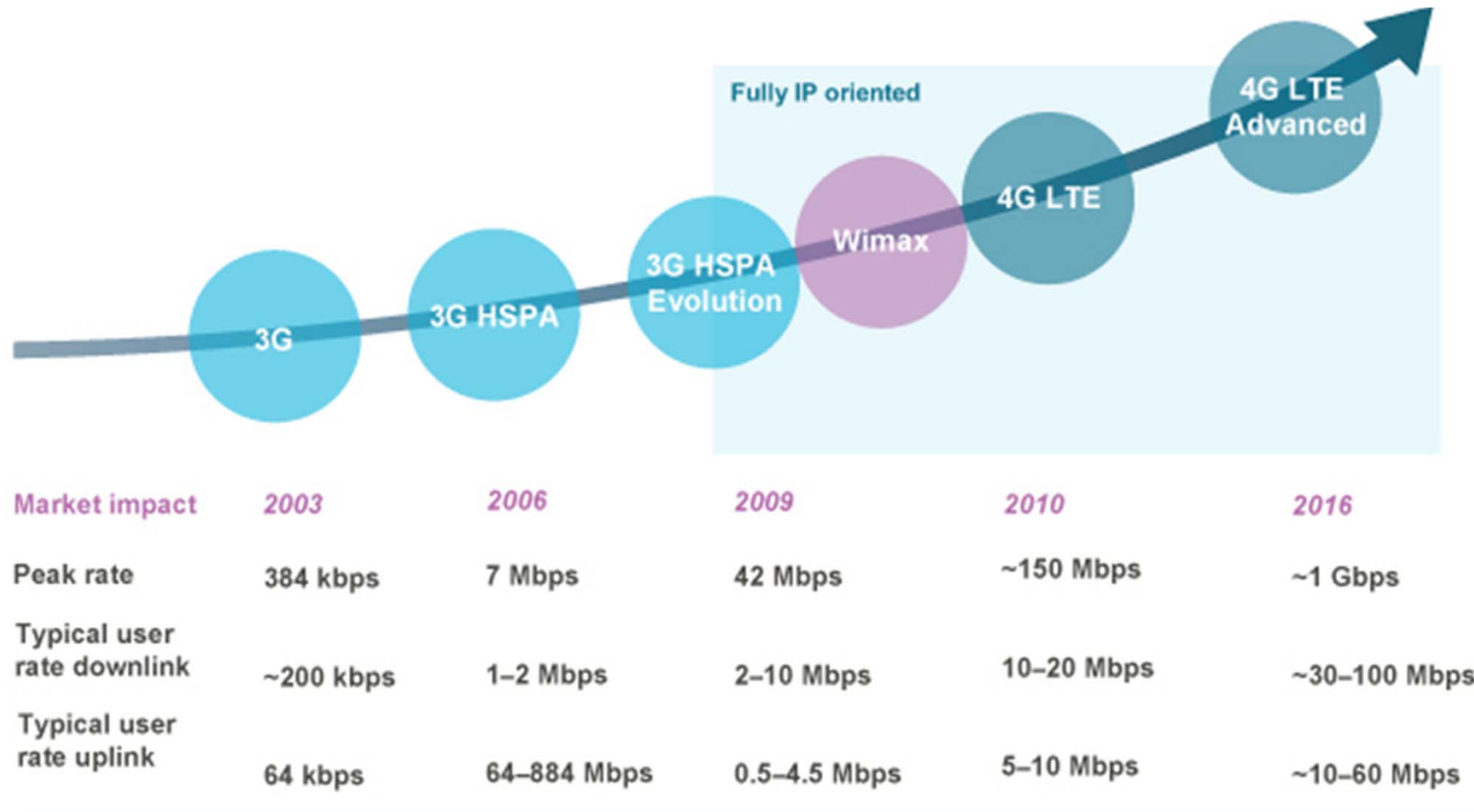
| 2014      | 2019      |
|-----------|-----------|
| 20.3 Mbps | 42.5 Mbps |

More Video  
Viewing



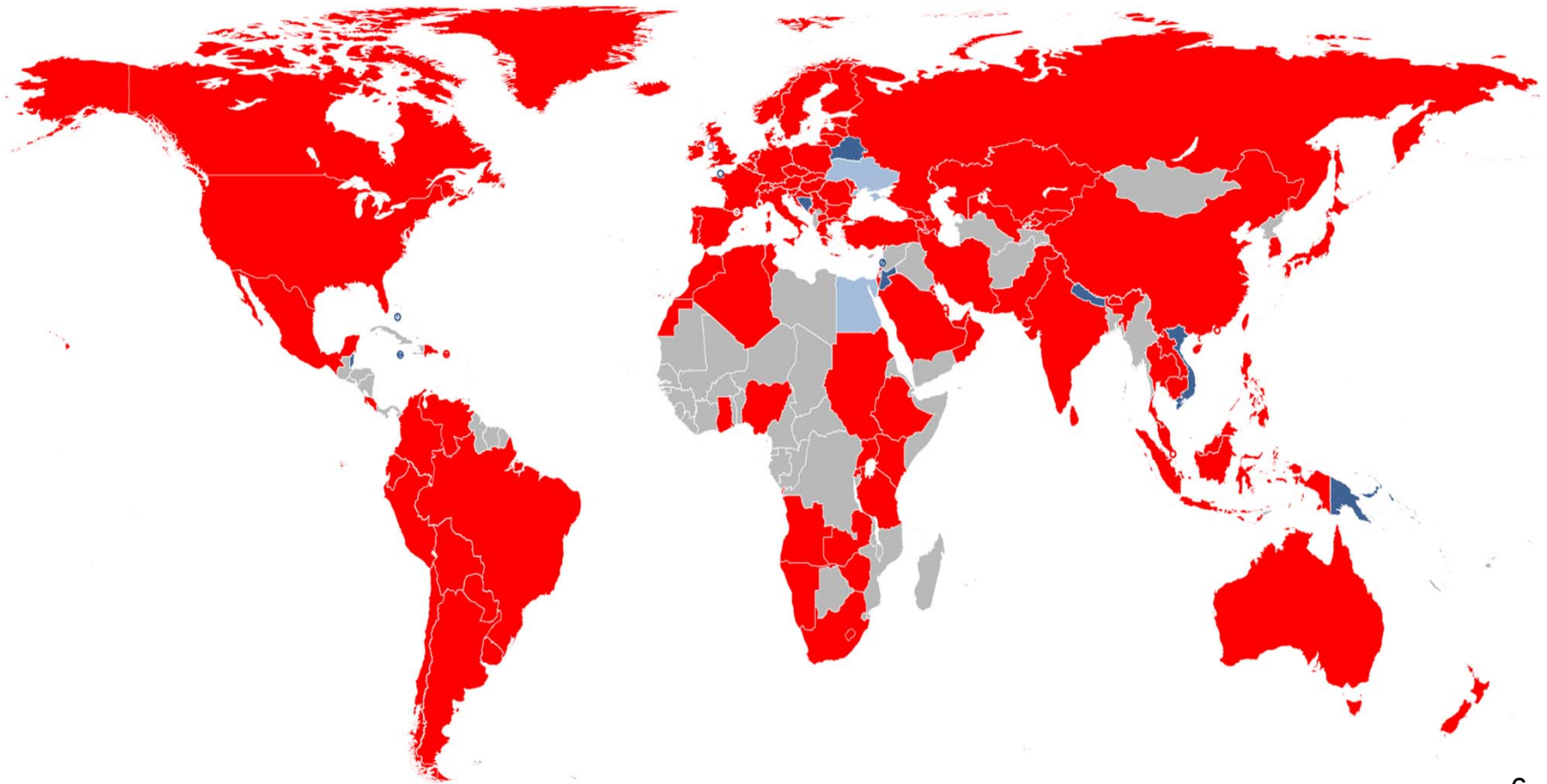
| 2014              | 2019              |
|-------------------|-------------------|
| 67% of<br>Traffic | 80% of<br>Traffic |

# Evolution of cellular communication



# Support of LTE-A in the world

- Red: in operation; blue: in planning

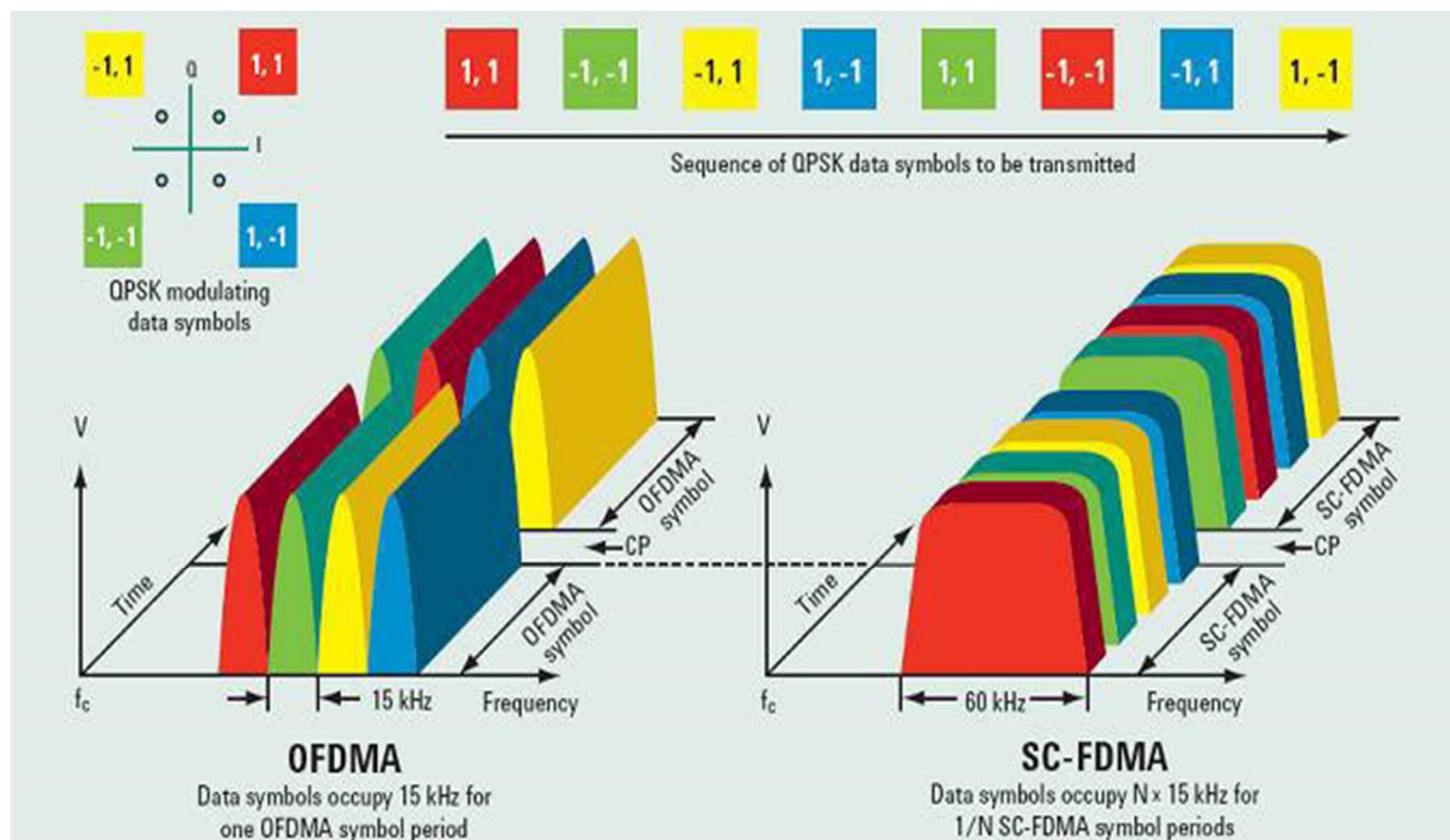


# Transmission technology in LTE-A

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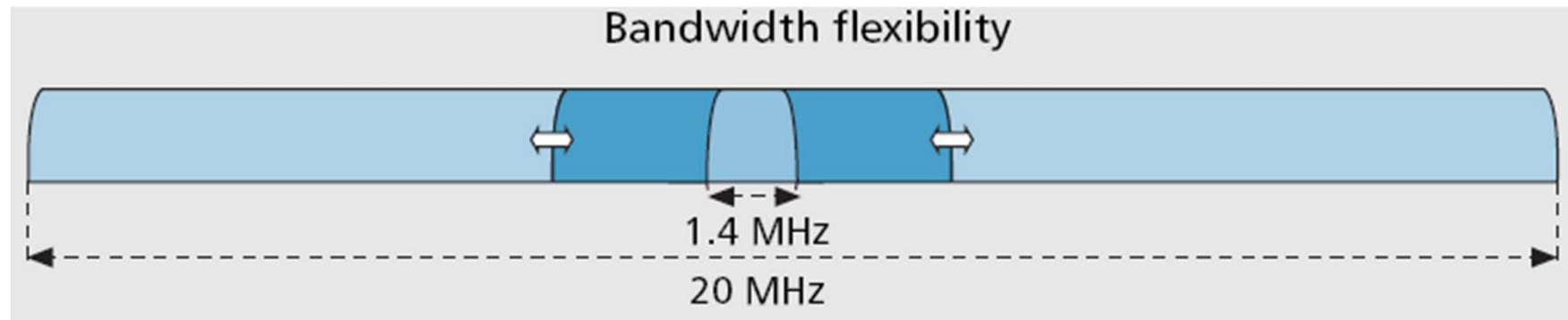
- LTE-A uses **OFDMA** and **SC-FDMA** for downlink and uplink communication, respectively.
  - OFDMA: Orthogonal frequency division multiple access
  - SC-FDMA: Single carrier FDMA
- Comparing with OFDMA, SC-FDMA has a lower PAP (peak to average power) ratio, which can save the *energy consumption* of mobile phones.

# OFDMA vs. SC-FDMA



# LTE-A channel bandwidth

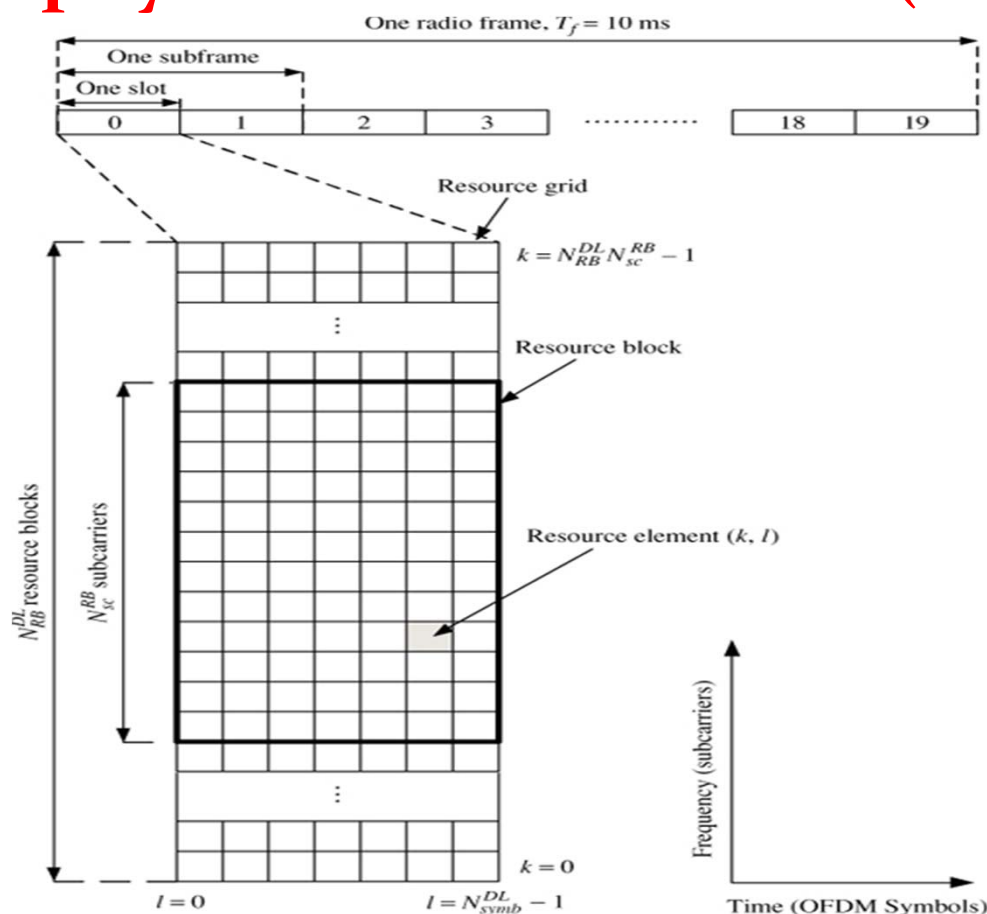
- LTE-A supports six types of channel bandwidth, including 1.4, 3, 5, 10, 15, and 20MHz.



- With *carrier aggregation*, a base station can combine different channels to provide much larger bandwidth (to **100MHz**).

# LTE-A resource

- The basic unit of LTE-A resource is called **physical resource block (PRB)**.



One PRB  
Time: 0.5ms  
Frequency: 12 subcarriers

# Outline

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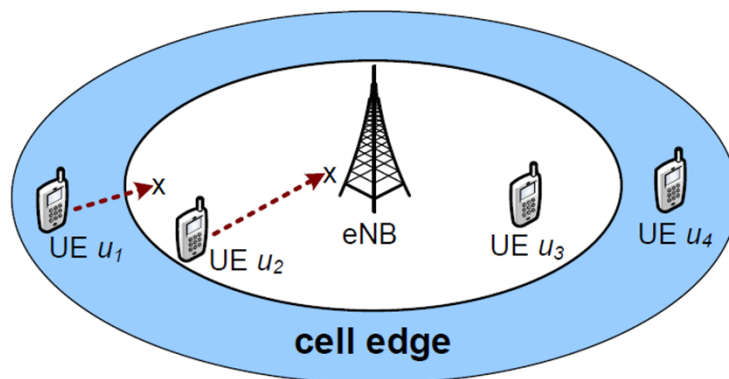
- Brief introduction to LTE-A
- A 4-guideline downlink scheduling strategy
- Downlink resource management for green communication

# Related work

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- Many downlink scheduling methods are proposed to allocate PRBs.
  - Max-CQI: (CQI – channel quality indicator)
    - Improve network throughput, but could *starve* UEs with bad channel quality.
  - Proportional Fair (PF):
    - Keep fairness by considering *current rate*  $r_i$  and *average rate*  $r_i^{\text{avg}}$  (i.e., pick UE with  $\max r_i / r_i^{\text{avg}}$ ).
      - Modified Largest Weighted Delay First (M-LWDF)
      - Exponential Proportional Fair (EXP/PF)

# Observations



- Unfair transmission by PF (for **cell-edge** UEs):
  - Two UEs move and  $r_1 / r_1^{\text{avg}} < r_2 / r_2^{\text{avg}}$ , so  $u_2$  is preferred by PF (but  $u_1$  requires more resource).
  - Two UEs stay and  $r_3 / r_3^{\text{avg}} \approx r_4 / r_4^{\text{avg}}$ , so  $u_3$  and  $u_4$  are treated as *no difference* (but  $u_4$  requires more resource).

# Network model (1/2)

- We consider SU-MIMO, so PRBs are *non-sharable*.
- LTE supports three modulations:
  - **QPSK** (CQI: 1~6), **16QAM** (CQI: 7~9), and **64QAM** (CQI: 10~15).

| CQI index | modulation   | code rate x 1024 | efficiency |
|-----------|--------------|------------------|------------|
| 0         | out of range |                  |            |
| 1         | QPSK         | 78               | 0.1523     |
| 2         | QPSK         | 120              | 0.2344     |
| 3         | QPSK         | 193              | 0.3770     |
| 4         | QPSK         | 308              | 0.6016     |
| 5         | QPSK         | 449              | 0.8770     |
| 6         | QPSK         | 602              | 1.1758     |
| 7         | 16QAM        | 378              | 1.4766     |
| 8         | 16QAM        | 490              | 1.9141     |
| 9         | 16QAM        | 616              | 2.4063     |
| 10        | 64QAM        | 466              | 2.7305     |
| 11        | 64QAM        | 567              | 3.3223     |
| 12        | 64QAM        | 666              | 3.9023     |
| 13        | 64QAM        | 772              | 4.5234     |
| 14        | 64QAM        | 873              | 5.1152     |
| 15        | 64QAM        | 948              | 5.5547     |

## Network model (2/2)

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- Given CQI and demand of each UE, our problem asks how to allocate PRBs to UEs such that
  - Network *throughput* is increased.
  - System *fairness* is improved.
  - Real-time packet *dropping* is reduced.

# 4G-LDS strategy

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- We develop a *4-guideline LTE downlink scheduling (4G-LDS) strategy* to solve the problem.
  - ① Credit-based guideline
  - ② Cell-division guideline
  - ③ Flow-weight guideline
  - ④ Packet-fitness guideline

# Credit-based guideline (1/2)

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- Motivated by *weighted fair queuing*, we expect that each UE can get  $t_i$  resource in every TTI.
  - If all UEs each exactly receives  $kt_i$  amount of data in  $k$  TTIs, for any  $k$ , we say that the network is **fair**.
- We use  $D_i$  to record the difference between *data actually received* by a UE  $u_i$  and *resource that it is expected to obtain*.
  - $D_i > 0$ :  $u_i$  uses more resource than expectation.
  - $D_i < 0$ :  $u_i$  does not receive enough resource.

## Credit-based guideline (2/2)

- The variation of  $D_i$  may be large, especially when some UEs have *much better channel quality* but others do not.

- We thus translate  $D_i$  into a **normalized** credit:

$$\hat{C}_i = 2 - \frac{D_i - D_{\min}}{D_{\max} - D_{\min}}$$

- This normalized credit will be limited to  $[1, 2]$ .

# Cell-division guideline (1/2)

- Cell-edge UEs do not have good channel quality and may be **starved** (e.g., max-CQI).
- We thus reserve a number of PRBs for them:

$$\gamma = \left\lceil \min \left\{ \frac{\sum_{u_i \in \mathcal{U}_e} \alpha_i}{\sum_{u_i \in \mathcal{U}} \alpha_i}, \beta \right\} \times m \right\rceil$$

- $\mathcal{U}$  and  $\mathcal{U}_e$  are sets of *all* UEs and *cell-edge* UEs.
- $\alpha_i$  is the demand of a UE.
- $\beta$  avoids keeping too many PRBs for cell-edge UEs.
- $m$  is the number of available PRBs.

# Cell-division guideline (2/2)

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- How to find cell-edge UEs?
  - One solution is to use the *distance*, but it requires the knowledge of UEs' locations.
- We propose a simple solution by **CQI**.
  - A UE is viewed as a cell-edge UE if it has  $CQI \leq 6$ .
  - This solution adds *no overhead*, since every UE has to report its CQI in each TTI.
  - When  $CQI \leq 6$ , the UE can use only the simplest **QPSK** modulation.

# Flow-weight guideline

- We assign different **weights** for real-time and non-real-time flows.

- Real-time flows: (modified *M-LWDF*)

$$\hat{W}_{i,j} = w_{i,j} \times d_{i,j} \times r_i$$

- $d_{i,j}$ : HOL packet delay of the flow
- $r_i$ : data rate of the UE

- Non-real-time flows: (modified *max-CQI*)

$$\hat{W}_{i,j} = r_i$$

# Packet-fitness guideline (1/2)

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- We use a *fitness degree* to check if the PRB is **suitable** to send the flow's packet.
  - $l_{i,j}$  is the size of the flow's HOL packet.
  - $\theta_i$  is the PRB's capacity (i.e., # of bits carried).
  - $\Theta_{64\text{QAM}}$  and  $\Theta_{16\text{QAM}}$  are the minimum capacity of an PRB when 64QAM and 16QAM are adopted.
- Three cases are discussed below.

# Packet-fitness guideline (2/2)

- $\theta_i \geq \Theta_{64\text{QAM}}$  (large-capacity PRB):
  - Prefer matching the PRB with a *large-size packet*.
  - Fitness degree:  $\hat{F}_{i,j} = \min\{l_{i,j}/\theta_i, 1\}$
- $\theta_i$  is in  $[\Theta_{16\text{QAM}}, \Theta_{64\text{QAM}})$ :
  - Fitness degree =  $\theta_i / \Theta_{64\text{QAM}}$  if  $l_{i,j} \geq \Theta_{64\text{QAM}}$ .
  - Otherwise,  $\hat{F}_{i,j} = \min\{l_{i,j}/\theta_i, 1\}$
- $\theta_i < \Theta_{16\text{QAM}}$ :
  - Fitness degree =  $\theta_i / \Theta_{16\text{QAM}}$  if  $l_{i,j} \geq \Theta_{16\text{QAM}}$ .
  - Otherwise,  $\hat{F}_{i,j} = \min\{l_{i,j}/\theta_i, 1\}$

# Discussion

- Our 4G-LDS strategy will conduct these steps:
  - Discard those **overdue** packets (i.e., cannot meet their deadlines even if we *send them out now*) in advance to save bandwidth.
  - Reserve  $\gamma$  PRBs for cell-edge UEs, and allow **all** UEs to compete for the remaining  $(m - \gamma)$  PRBs.
  - Select a flow to receive each PRB by

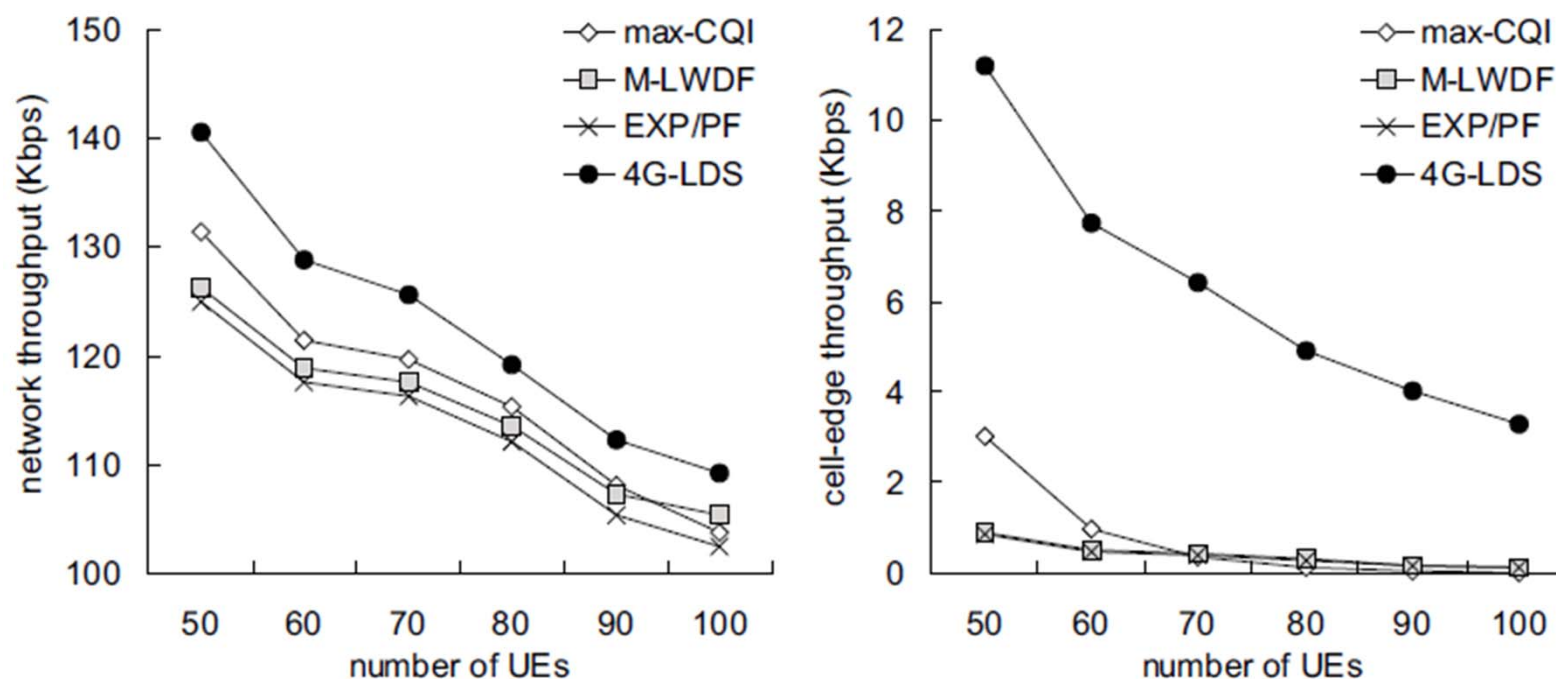
$$f_{i,j} = \arg \max_{i,j} \begin{cases} \hat{C}_i \times \hat{F}_{i,j} \times r_i & \text{if } u_i \in \mathcal{U}_e, \\ \hat{C}_i \times \hat{F}_{i,j} \times \hat{W}_{i,j} & \text{otherwise.} \end{cases}$$

# Simulation parameters

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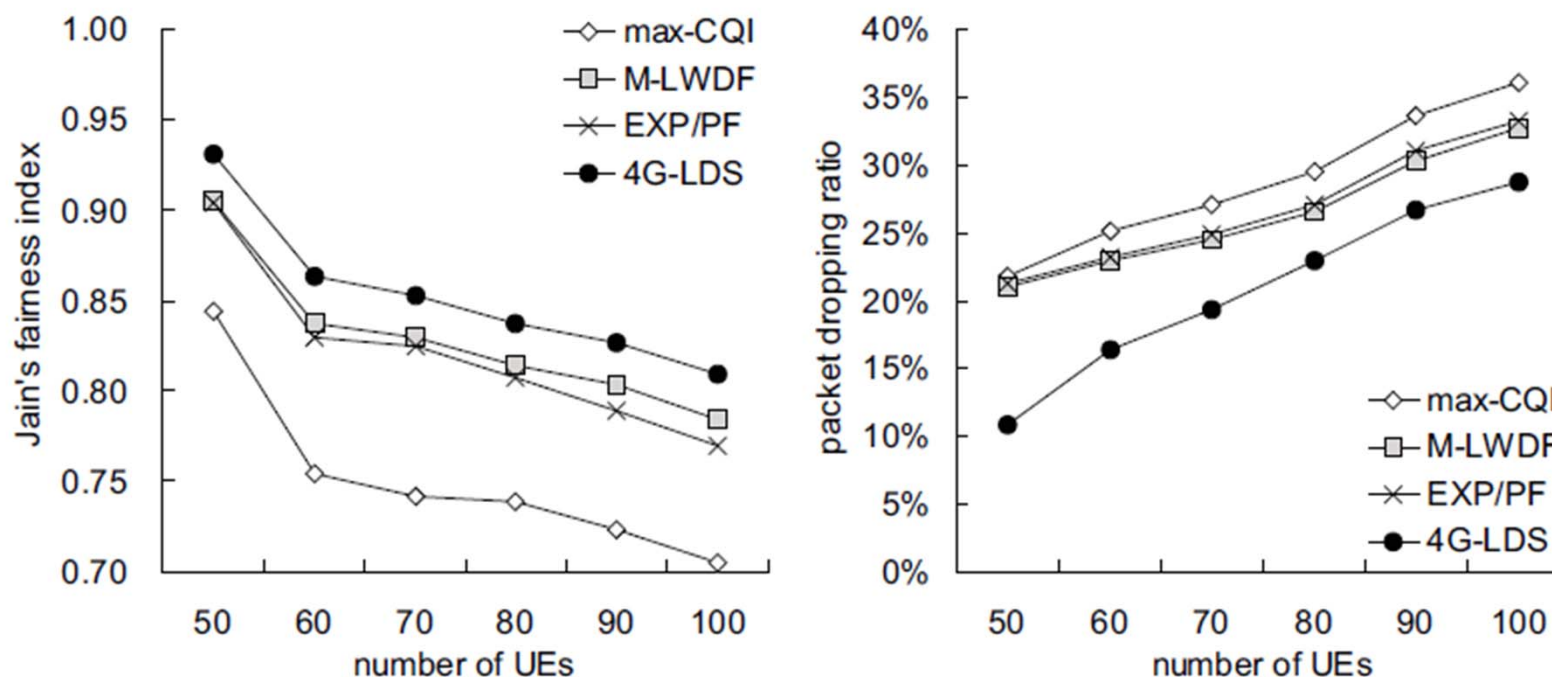
- We use **LTE-Sim** to simulate an LTE macro-cell.
  - A 20MHz downlink channel with 100 PRBs
  - 50 ~ 100 UEs in the cell
  - Mobility model:  
random walking with 3km/h speed
  - Each UE has two flows:
    - Video streaming (242Kbps)
    - Constant-bit-rate (CBR, 12Kbps)
  - Comparison: max-CQI, M-LWDF, and EXP/PF

# Simulation results (1/2)



- 4G-LDS improves network throughput, especially for those *cell-edge UEs*.

# Simulation results (2/2)



- 4G-LDS not only improves *fairness*, but also reduces real-time *packet dropping*.

# Outline

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# Carrier aggregation (CA)

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- LTE-A supports channel bandwidth to 100MHz.
  - However, many bands in **microwave** spectrum have been used by 2G and 3G systems.
  - LTE-A uses CA to combine different *component carriers* (CCs) to get larger bandwidth.
    - Ex: Combine five 20MHz CCs to get 100MHz bandwidth.
- LTE-A allows *contiguous* and *non-contiguous* CA.
  - However, user equipment will consume more energy when non-contiguous CA is used.

# Related work

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- Existing methods are classified into two groups.
  - CC selection:
    - How to *balance the loads* of different CCs to improve the overall throughput?
  - PRB assignment:
    - How to allocate PRBs to each user equipment (UE) to improve throughput, or *reduce packet delay*?
- However, the issue of saving **energy** on communication is not well addressed.

# Problem formulation

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- Energy-reduction LTE-A resource management (ELRM) problem asks to allocate RBs such that
  - *Energy consumption* of each UE is minimized.
  - BS's *transmission power* is minimized.
- ELRM has two constraints:
  - Total power on all CCs cannot exceed the BS's *maximum power*  $P_{\max}$ .
  - *Traffic demand* of each UE should be met.
- ELRM is **NP-hard**.

# The proposed solution (1/4)

- Given a set  $\mathbf{U}$  of UEs (each with demand  $r_i$ ) and a set  $\mathbf{C}$  of CCs, our solution has the following steps.
  - ① Set an *initial power* of  $P_{\max}/|\mathbf{C}|$  on each CC.
  - ② BS then sends a *reference signal* to let each UE compute its **CQI** (channel quality indicator).
  - ③ With CQI, BS can find **MCS** (modulation & coding scheme) for each (UE, CC) pair.
- It thus builds a table where each tuple  $(u_i, c_j)$  is *the number of bits* sent by CC  $c_j$  for UE  $u_i$ .

## The proposed solution (2/4)

- ④ We also use  $Q_i$  to record the **backlog** of each UE (i.e., the number of bits *not satisfied yet*).
- ⑤ BS picks a UE with maximum  $Q_i$  and assigns one CC by its **weight**.
  - Weight is the *sum of data bits* supported by  $c_j$  for all UEs (except  $u_i$ ) with positive backlog.
  - A large weight indicates that other UEs may have *good channel quality* on  $c_j$ , so we should keep it for later assignment.

# The proposed solution (3/4)

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- ⑤ (Cont.) BS will select the CC with *largest data rate* but *smallest weight* to serve the UE.
  - We then update  $Q_i$  (backlog) and corresponding tuples in the table.
- ⑥ We repeat step 5 until any of two cases occurs:
  - The demand of each UE is satisfied.
    - Go to step 7 to save *BS's transmission power*.
  - All CCs are used but some UEs have backlog.
    - Algorithm **finishes** as there is no enough resource.

# The proposed solution (4/4)

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- ⑦ We check if some CCs can be given with less **power** to meet UEs' demands.
- Suppose that a CC  $c_j$  is assigned to a UE  $u_i$  with MCS  $m(i, j)$ .
  - We try a new MCS  $m'(i, j) < m(i, j)$  and check if  $c_j$  has *enough PRBs* to satisfy  $u_i$ 's demand.
  - If so, we can lower down the power on  $c_j$  such that the MCS becomes  $m'(i, j)$ .

## Example (1/3)

| UE     | Demand | $c_1$ | $c_2$ | $c_3$ | $c_4$ | $c_5$ |
|--------|--------|-------|-------|-------|-------|-------|
| $u_1$  | 1150   | 380   | 884   | 0     | 380   | 380   |
| $u_2$  | 550    | 236   | 380   | 0     | 884   | 153   |
| $u_3$  | 594    | 380   | 606   | 0     | 236   | 606   |
| weight |        |       |       |       |       |       |

- ① Pick  $u_3$  and CC  $c_2$ ,  $c_3$ , and  $c_5$  are candidates.
  - Select  $c_3$  and update  $Q_3$  by  $1200 - 606 = 594$ .
- ② Pick  $u_1$  and CC  $c_2$  is the only candidate.
  - Assign  $c_2$  to  $u_1$  and update  $Q_3$  by  $1150 - 884 = 266$ .

## Example (2/3)

| UE    | Demand | $c_1$ | $c_2$ | $c_3$ | $c_4$ | $c_5$ |
|-------|--------|-------|-------|-------|-------|-------|
| $u_1$ | 266    | 380   | 0     | 0     | 126   | 0     |
| $u_2$ | 0      | 236   | 0     | 0     | 294   | 0     |
| $u_3$ | 0      | 380   | 0     | 0     | 78    | 0     |

- ③ Pick  $u_3$  again and CC  $c_5$  is the candidate.
  - Assign  $c_5$  and update  $Q_3$  by zero.
- ④ Pick  $u_2$  and CC  $c_4$  is the only candidate.
  - UE  $u_2$  only uses 8 PRBs in CC  $c_4$ .

## Example (3/3)

| UE    | Demand | $c_1$ | $c_2$ | $c_3$ | $c_4$ | $c_5$ |
|-------|--------|-------|-------|-------|-------|-------|
| $u_1$ | 0      | 0     | 0     | 0     | 126   | 0     |
| $u_2$ | 0      | 0     | 0     | 0     | 294   | 0     |
| $u_3$ | 0      | 0     | 0     | 0     | 78    | 0     |

⑤ Assign  $c_1$  to  $u_1$ .

- CC assignment:

$\{c_1, c_2\} \rightarrow u_1$

$\{c_4\} \rightarrow u_2$  **We can lower down MCS in  $c_4$ .**

$\{c_3, c_5\} \rightarrow u_3$

# Simulation parameters

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- Path loss:  $PL = 128.1 + 37.6 \log D(\text{BS}, u_i)$
- Two communication bands:
  - Band 1 (2110~2170 MHz)  
→ cut in twelve 5MHz CCs
  - Band 5 (869~894 MHz)  
→ cut in two 5MHz CCs and five 3MHz CCs
- Each UE has 50Mbps real-time traffic and 35Mbps non-real-time traffic.
- BS's maximum power: 40W

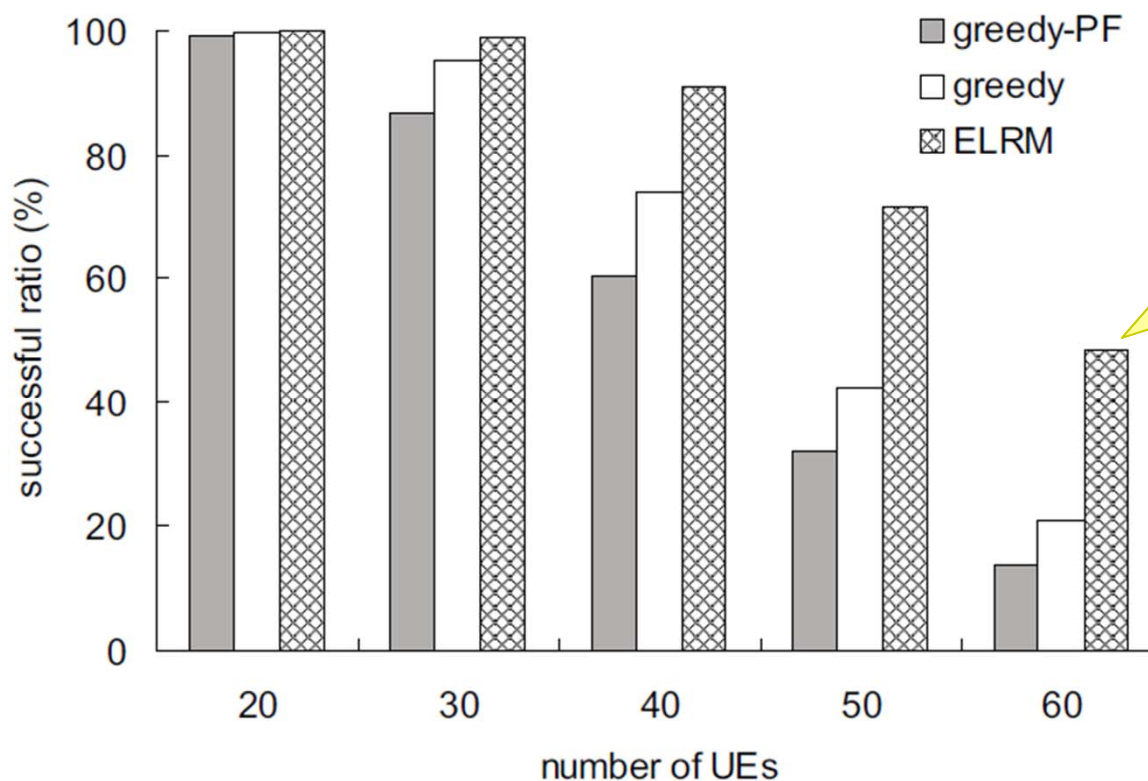
# Comparison

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- Two methods for comparison:
  - Greedy-PF: Find a **weighted** rate of each UE and selects the pair of UE and CC with the maximum rate. To keep fairness, a larger weight is given to a UE that *sent less data* in the past.
  - Greedy: Similar with greedy-PF but set all weights to one.
- For each experiment, we repeat 1000 simulations and take their average.

# Experimental results (1/3)

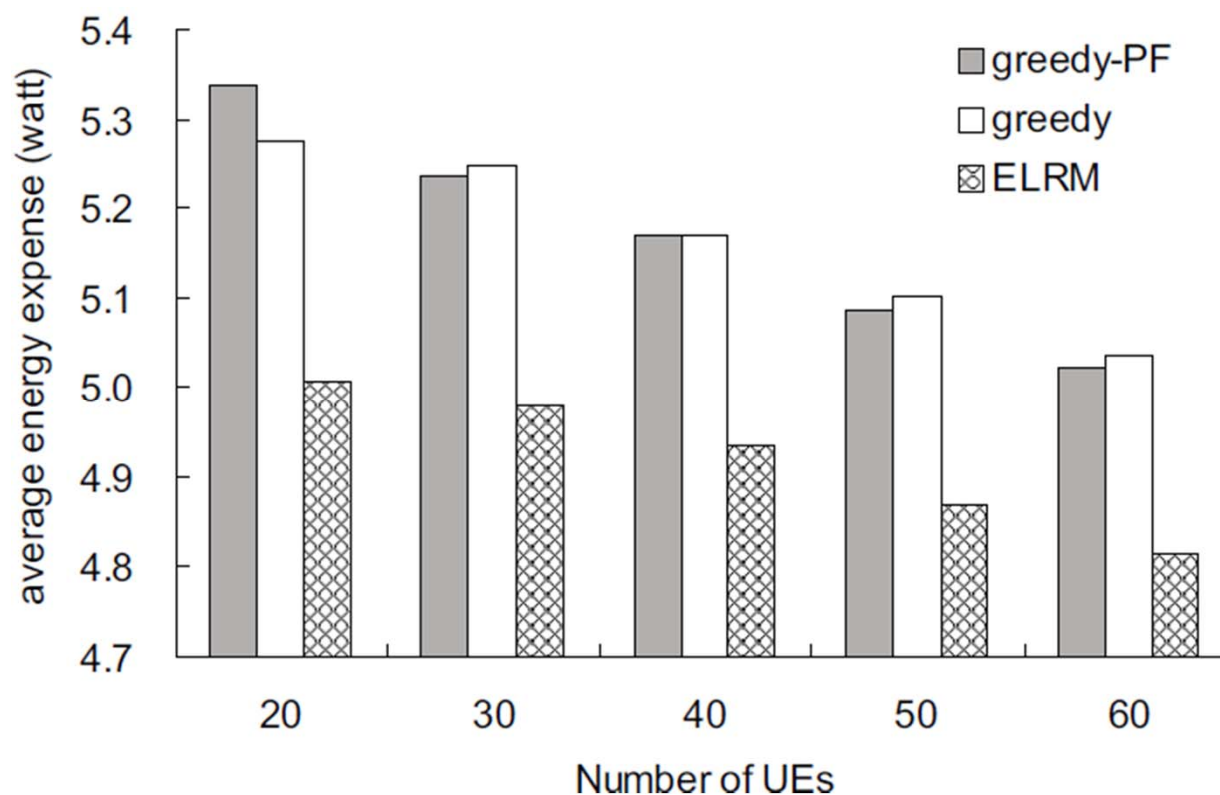
- Successful ratio of resource scheduling:
  - # of successful experiments / 1,000 experiments



Our solution can still keep 50% of successful ratio.

## Experimental results (2/3)

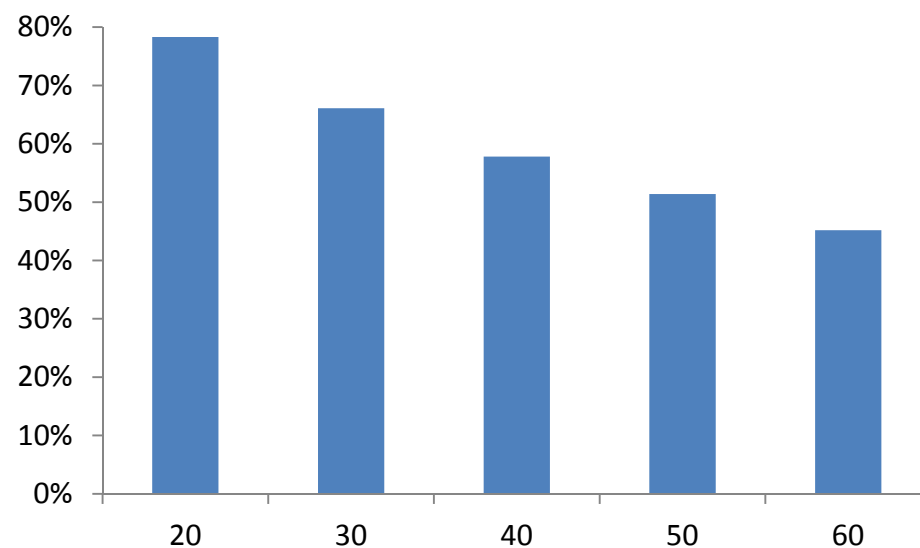
- Average energy expense of UEs:
  - Decrease as the number of UEs grows.



## Experimental results (3/3)

- Saving ratio of the BS's transmission power by our ELRM solution:

| UEs   | 20    | 30    | 40    | 50    | 60    |
|-------|-------|-------|-------|-------|-------|
| ratio | 78.3% | 66.1% | 57.8% | 51.4% | 45.2% |



# Conclusion

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- LTE-A is a world-wide standard for 4G cellular communication.
- We have developed two downlink resource scheduling methods for LTE-A networks.
  - 4-guideline strategy
  - Green communication

# References

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- You-Chiun Wang\* and Dai-Rong Jhong, "Efficient allocation of LTE downlink spectral resource to improve fairness and throughput", *International Journal of Communication Systems*, accepted, 2017.
- You-Chiun Wang\* and Hung-Yi Ko, "Energy-efficient downlink resource scheduling for LTE-A networks with carrier aggregation," *Journal of Information Science and Engineering*, vol. 33, no. 1, pp. 123-141, Jan. 2017.